



# Chapter5 Logical Database Design and the Relational Model

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Logical database design and  
the relational model



# Outline

- ◆ **Introduction**
- ◆ The relational data model
- ◆ Transforming EER diagrams into relations
- ◆ Normalization, Normal forms, relational schema decomposition
- ◆ Merging relations

# Introduction

## ◆ **Logical database design**

the conceptual data model →

relational schemas as the logical data model

- relational model

- E-R and EER model → relational model

The objective of logical database design:

translate the conceptual design into a logical database design that can be implemented on a chosen DBMS.

# The relational data model

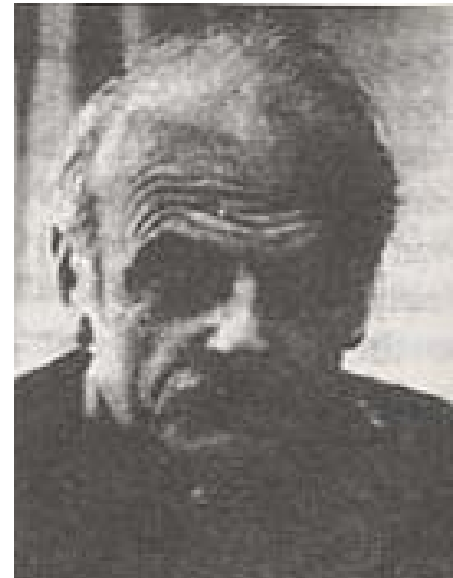
- ◆ **History**

- Relational data model (E .F. Codd, 1970)
- System R
- Ingres
- Commercial RDBMS (1980)

- ◆ **Basic Definitions**

- ◆ **Integrity constraints**

- ◆ **Creating relational tables and well-structured relations**



# Basic definitions (1)

## ◆ 3 components of the relational data model

- Data structure: tables
- Data manipulation: operations (SQL)
- Data integrity: constraints



## ◆ Relational data structure

- Relation: a named, 2-dimensional table of data
- Relation — table; Row — tuple, record; Column — attribute
- E.g. Employee1(Emp\_id, Name, Dept\_Name, Salary)

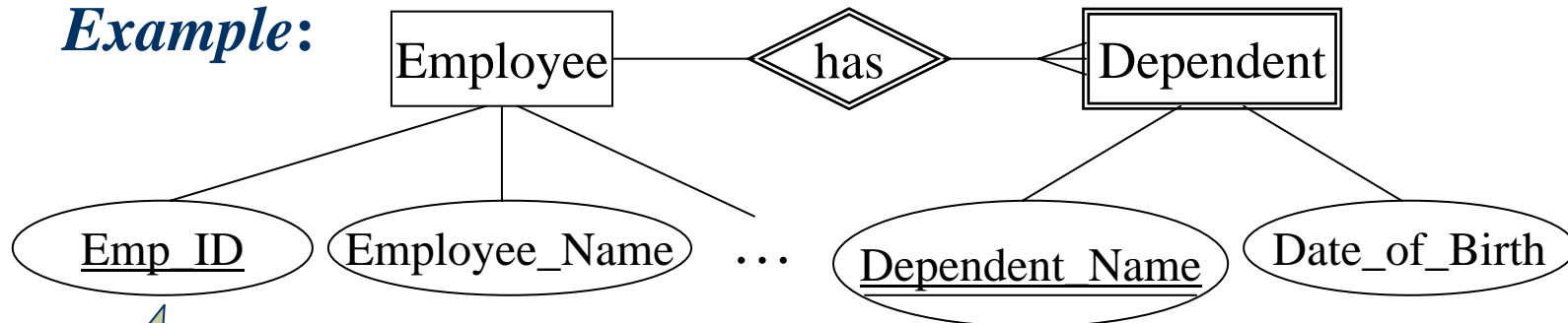
<u>Emp_id</u>	Name	Dep_Name	Salary
100	Margaret	Marketing	48,000
140	Allen	Accounting	52,000

# Basic definitions (2)

## ◆ Relational keys

- **Primary key (identifier)**: an attribute (or composition of attributes) that uniquely identifies each row in a relation — Entity integrity

*Example:*



Attribute as  
primary key

The identifier of “*Dependent*”:  
Composition of attributes as primary key  
**Dependent (Emp\_id, Dependent Name, Data\_of\_Birth)**

# Basic definitions (3)

## ◆ Relational keys

- **Foreign key**—Referential integrity

- ① An attribute in a relation of a database that serves as **the primary key attribute** of another relation in the same database
- ② Represents the relationship between two tables or relations
- ③ Imply the semantics of **contexts** and **relationships**

E.g., Department(Dept Name, Location, Fax)

Employee1(Emp\_id, Name, Dept\_Name, Salary)

Dependent (Emp\_id, Dependent Name, Data\_of\_Birth)

“Dept\_Name” is the foreign key of “Employee1”

“Emp\_id” is the foreign key of “Dependent”

# Basic definitions (4)

## ◆ Properties of relations

- Each relation (table) has a unique name
- No multivalued attributes in a relation (Atomic value)
- Each row is unique
- Each attribute has a unique name
- The sequence of columns and rows is insignificant

How to deal with the multivalued attributes?

Emp_id	Name	Dep_Name	Salary
100	Margaret	Marketing	48,000
140	Allen	Accounting	52,000

Emp_id	Name	Dep_Name	Salary
140	Allen	Accounting	52,000
100	Margaret	Marketing	48,000

Emp_id	Dep_Name	Name	Salary
100	Marketing	Margaret	48,000
140	Accounting	Allen	52,000



# An example

Customer (Customer\_id, Customer\_name, Address, ...)

Order (Order\_id, Customer\_id, Order\_date)

Product (Product\_id, Product\_description, Product\_finish, ...)

Order Line (Order\_id, Product\_id, Quantity)

Whether “Order” is a strong entity or weak entity?

How to insert data into above tables?

# Integrity constraints (1)

- ◆ **Domain constraints**

The set of values that may be assigned to an attribute:

domain name, meaning, data type, size, allowable values or range

- ◆ **Entity constraints**

- No primary key attribute (or component of a primary key attribute) can be null
- There do not exist any two tuples with the same primary key values

- ◆ **Referential constraints**

- A rule that states either each foreign key value must match a primary key value in another relation or the foreign key value must be null

# Integrity constraints (2)

## ◆ Referential constraints (Cont.)

- Identifying associations where referential integrity by means of the graphical representation using arrows in the schema
- Foreign key and cardinality

How do you know if a foreign key is allowed to be null?



Order (Order id, Customer\_id, Order\_date)

*Customer\_id* can not be NULL!



Student (Student id, Student\_name, Department\_id, ...)

*Department\_id* can be NULL!

# Integrity constraints (3)

## ◆ Referential constraints (Cont.)

- The foreign key attributes must match the existing primary key attributes
- The cardinality determines whether the foreign key can be null
- Whether a foreign key can be null must be specified as a property of the foreign key attribute when the database is defined
- The 3 choices when deleting the tuple that is referred due to foreign keys:
  - ① Cascading deletion
  - ② Prohibiting deletion
  - ③ Place a null value in the foreign key

How to carry them  
out in practice?

# Integrity constraints (4)

## ◆ Referential constraints (Cont.)

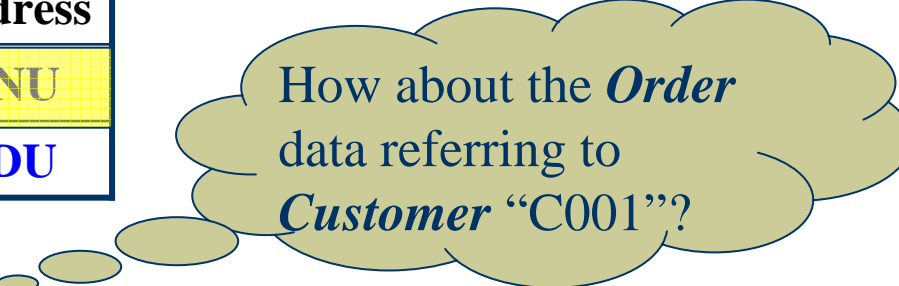
- Cascading deletion

Customer:

Customer_id	Customer_name	Address
C001	Mary	YNU
C002	John	FDU

Order:

Order_id	Customer_id	Order_date
O001	C001	2005-04-02
O002	C001	2005-04-05
O003	C002	2004-03-01



How about the *Order* data referring to *Customer* "C001"?

# Integrity constraints (5)

## ◆ Referential constraints (Cont.)

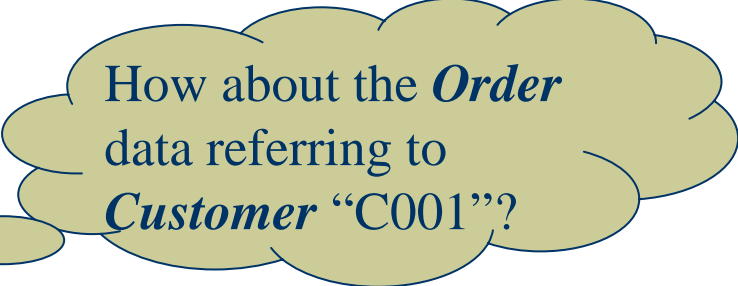
- Prohibiting deletion

Customer:

Customer_id	Customer_name	Address
C001	Mary	YNU
C002	John	FDU

Order:

Order_id	Customer_id	Order_date
O001	C001	2005-04-02
O002	C001	2005-04-05
O003	C002	2004-03-01



How about the *Order* data referring to *Customer* "C001"?

# Integrity constraints (6)

## ◆ Referential constraints (Cont.)

- Place a null value in the foreign key

Customer:

Customer_id	Customer_name	Address
C001	Mary	YNU
C002	John	FDU

Order:

Order_id	Customer_id	Order_date
O001	C001	2005-04-02
O002	C001	2005-04-05
O003	C002	2004-03-01

How about the *Order* data referring to *Customer* "C001"?

"Customer\_id" as the foreign key of *Order* can not be NULL!

# Well-structured relations (1)

- ◆ **Well-structured relations**

- Contains minimal redundancy
- Allows users to insert, modify and delete the rows in a table without errors or inconsistencies

- Why do these anomalies exist?

- How to get the well-structured relations?

- ◆ **Anomaly**

- **Insertion anomaly**: Null primary key value, nonexistent referred tuple, etc.
  - **Deletion anomaly**: The tuple in the parent table is deleted, but the tuples in the child table still exist, etc.
  - **Modification anomaly**: multiple times of modifications
- To the extreme, all the data can just stored in one relation!  
—— Universal relational schema (Discussion☺)



## Well-structured relations (2)

<u>Emp_id</u>	Name	Dept_name	Salary	<u>Course_title</u>	Date_completed
100	Simpson	Marketing	48,000	SPSS	6/19/200X
100	Simpson	Marketing	48,000	Surveys	10/7/200X
150	Martin	Marketing	42,000	Java	8/12/200X
...					

### Analysis:

Insertion anomaly, deletion anomaly; modification anomaly?



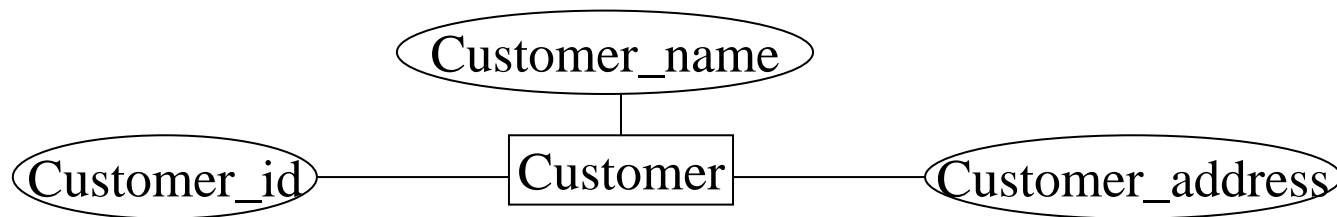
# Transforming EER diagrams into relations



- ◆ Step1: Map regular entities
- ◆ Step2: Map weak entities
- ◆ Step3: Map binary relationships
- ◆ Step4: Map associative entities
- ◆ Step5: Map unary relationships
- ◆ Step6: Map ternary (and n-ary) relationships
- ◆ Step7: Map supertype/subtype relationships

# Step 1: Map regular entities (1)

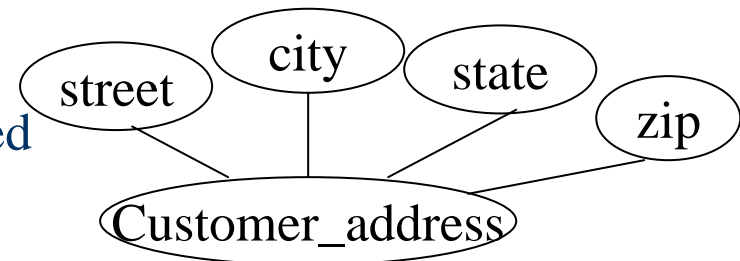
- ① - name of entity type → name of the relation  
- name of attribute of the entity → attribute of the relation



Customer	<u>Customer_id</u>	Customer_name	Customer_address
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- ② Composite attributes

- only the simple components are included



<u>Customer_id</u>	Customer_name	street	city	state	zip
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# Step1: Map regular entities (2)

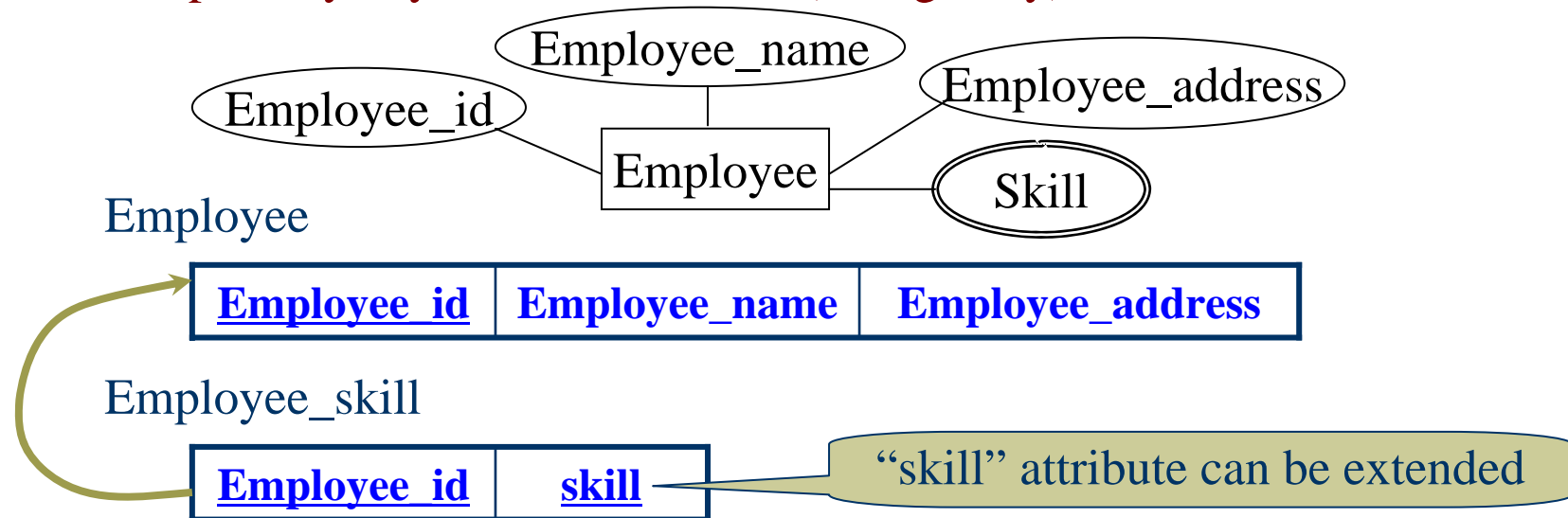
## ③ Multivalued attributes

2 relations:

1<sup>st</sup> relation: all the attributes except the multivalued attribute

2<sup>nd</sup> relation:

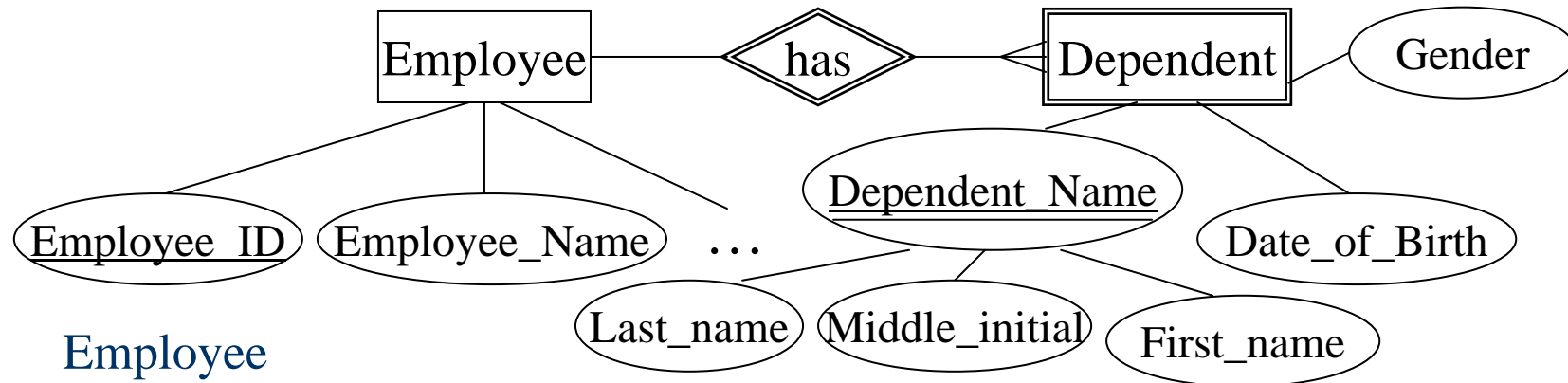
the primary key of the 1<sup>st</sup> relation (foreign key)+ multivalued attribute



# Step2: Map weak entities

Suppose the corresponding relation (strong entity) called identifying relation

- ① include all simple attributes of the weak entity
- ② **primary key = primary key of the identifying relation + partial identifier**



Employee

<u>Employee id</u>	Employee_name	Employee_address
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Dependent

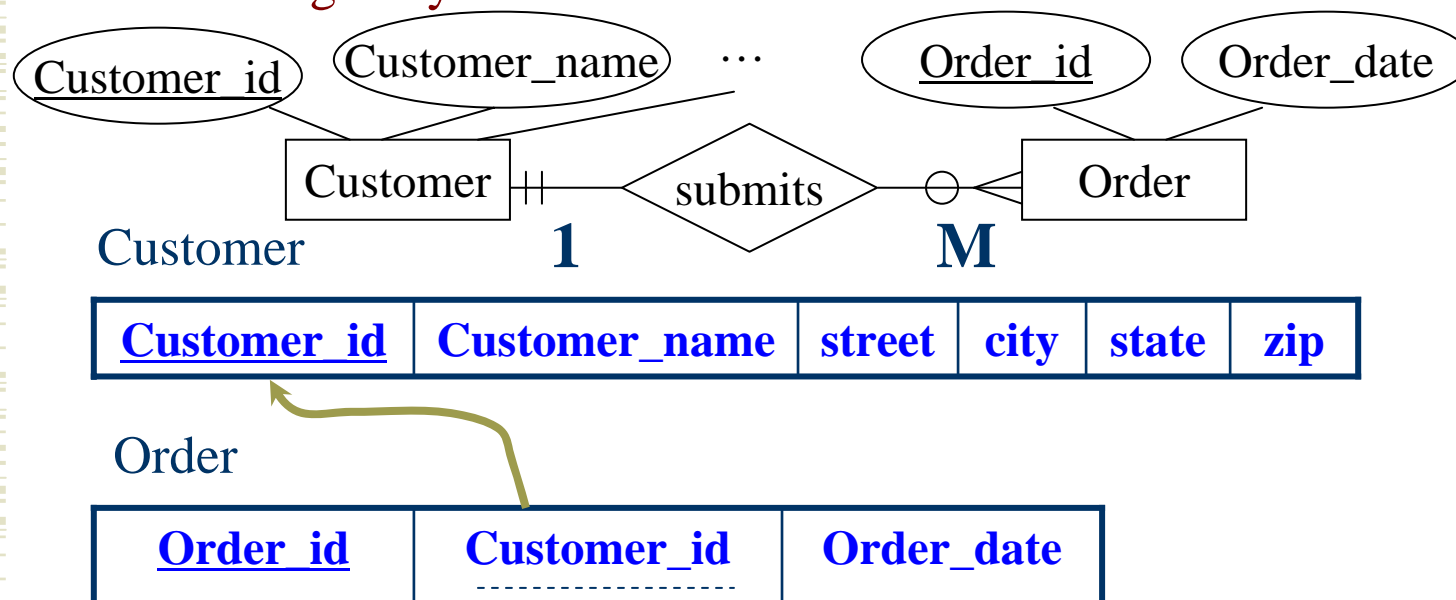
<u>Employee id</u>	<u>First name</u>	<u>Middle initial</u>	<u>Last name</u>	<u>Date of Birth</u>	Gender
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# Step3: Map binary relationships (1)

## ◆ Map binary one-to-many relationships (1:M)

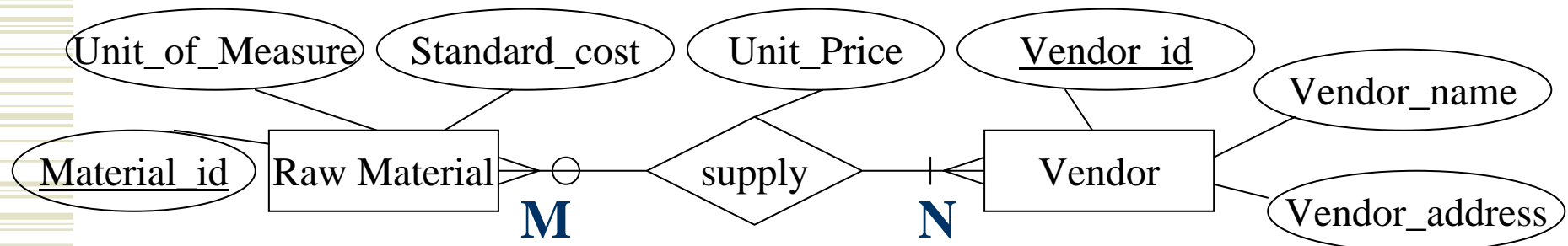
- ① create a relation for each of the two participating entity types
- ② add the primary key of **One** side relation to the **Many** side relation as the foreign key



# Step3: Map binary relationships (2)

- ◆ **Map binary many-to-many relationships (M:N)**

- ① create a relation for each of the two participating entity types: A, B
- ② create a new relation C, whose primary key is the combination of the primary keys of A and B, which are the foreign keys of C



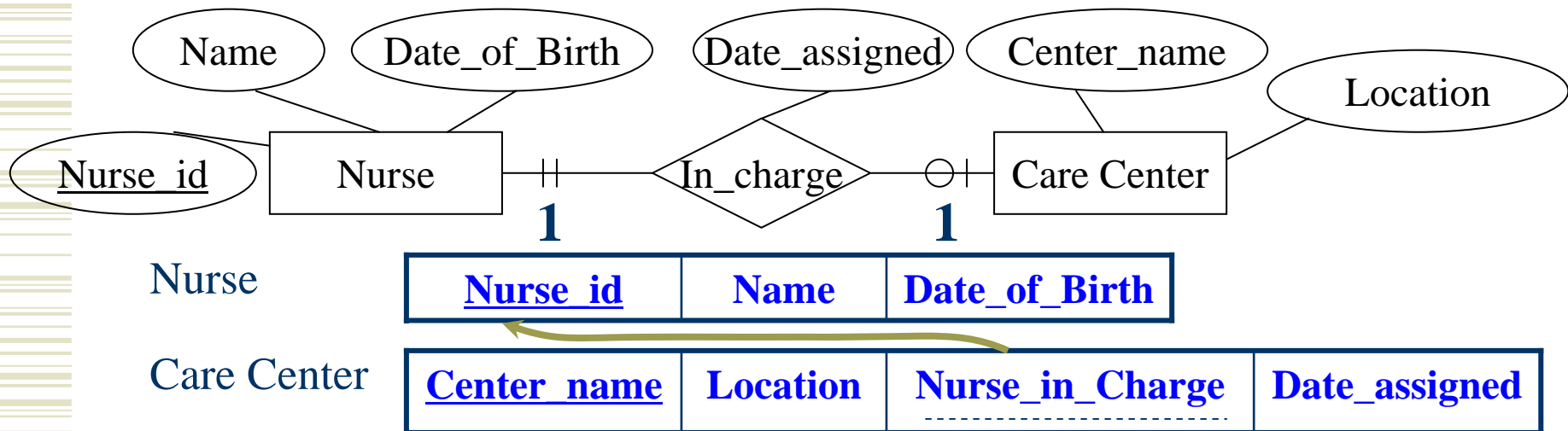
Raw Material	<u>Material id</u>	Standard_cost	Unit_of_Measure
Supply	<u>Material id</u>	<u>Vendor id</u>	Unit_Price
Vendor	<u>Vendor id</u>	Vendor_name	Vendor_address

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# Step3: Map binary relationships (3)

## ◆ Map binary one-to-one relationships (1:1)

- ① create a relation for each of the two participating entity types
- ② the primary key of one of the relations is included as a foreign key in the another relation



**Note:** Include the primary key of mandatory side to the optional side

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# Step4: Map associative entities (1)

Similar to step3

- one relation for each of the two participating entity types
- the third for the associative entity

Whether an identifier is assigned to the associative entity?

① Identifier not assigned:

The default primary key: the two primary key attributes from the other 2 relations, also as the foreign key respectively

② Identifier assigned:

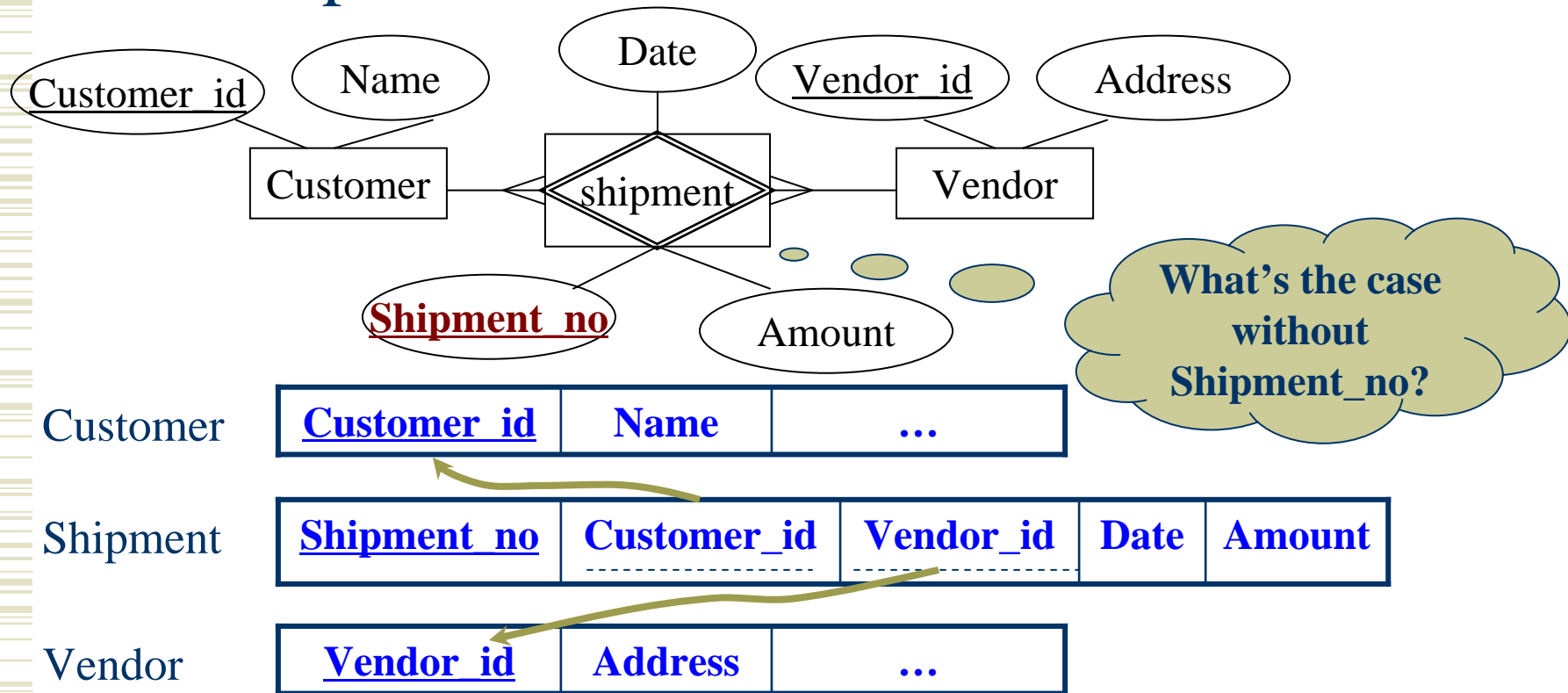
The default identifier may not uniquely identify instances of the associative entity

a) The identifier is the primary key of the associative entity

b) The primary keys for the 2 participating relations are the foreign keys in the associative relation

# Step4: Map associative entities (2)

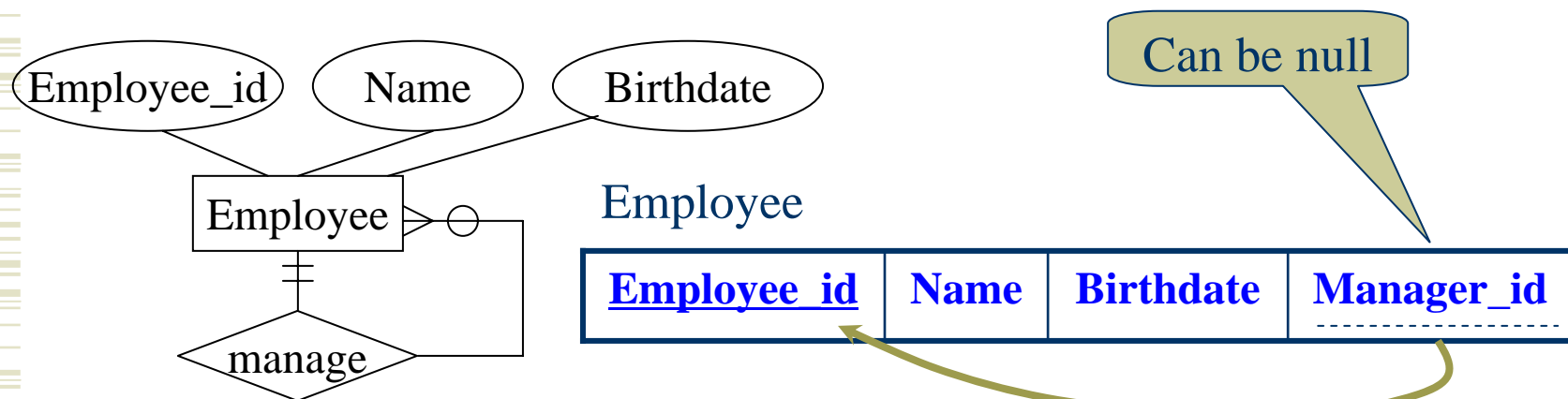
## ◆ Example



# Step5: Map unary relationship (1)

## ◆ Map unary one-to-many relationships:

- ① map the entity type into relation
- ② a foreign key attribute is added within the **same** relation that references the primary key values
  - the same domain as the primary key
  - recursive foreign key



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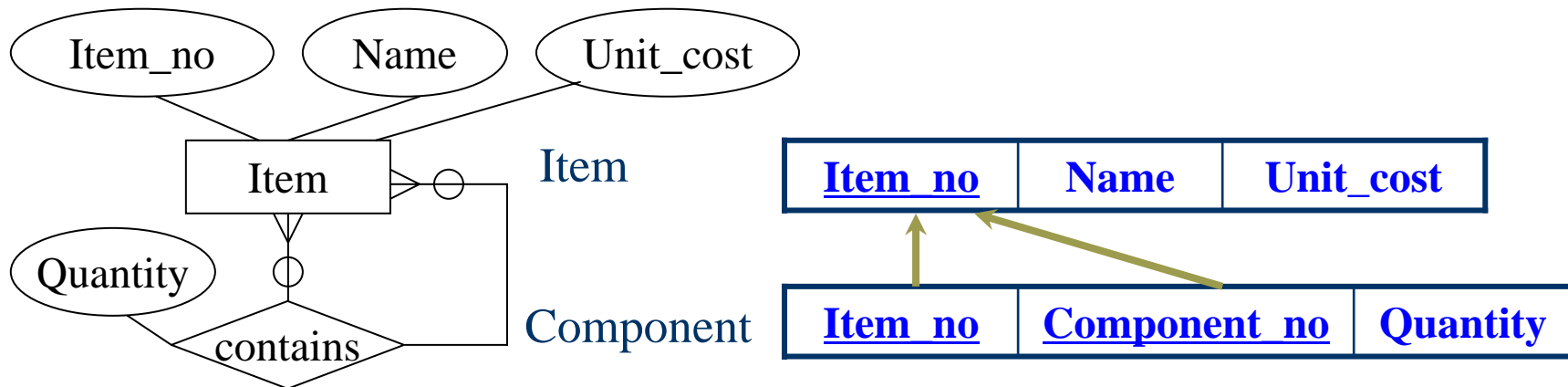
# Step5: Map unary relationship (2)

- ◆ **Map unary many-to-many relationships (BOM):**

2 relations:

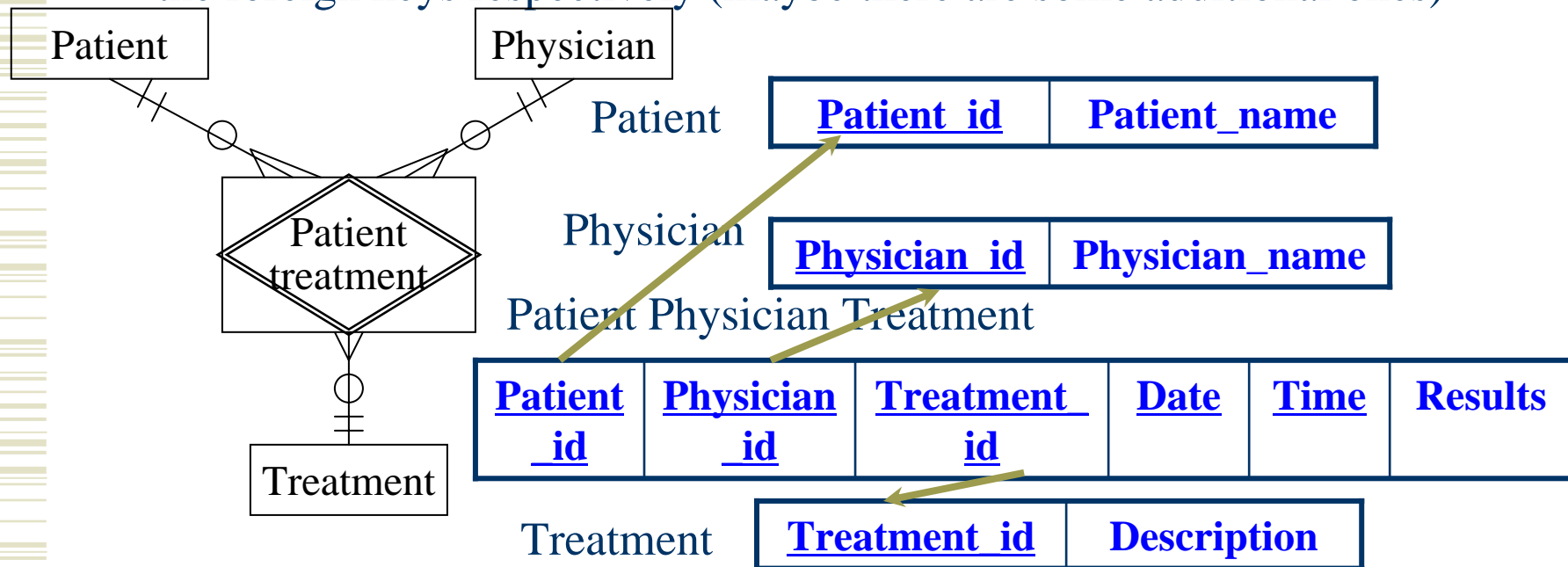
- ① one to represent the entity type in the relationship
- ② the other is an associative relation to represent the M:N relationship itself;

Primary key includes 2 attributes



# Step6: Map ternary relationships

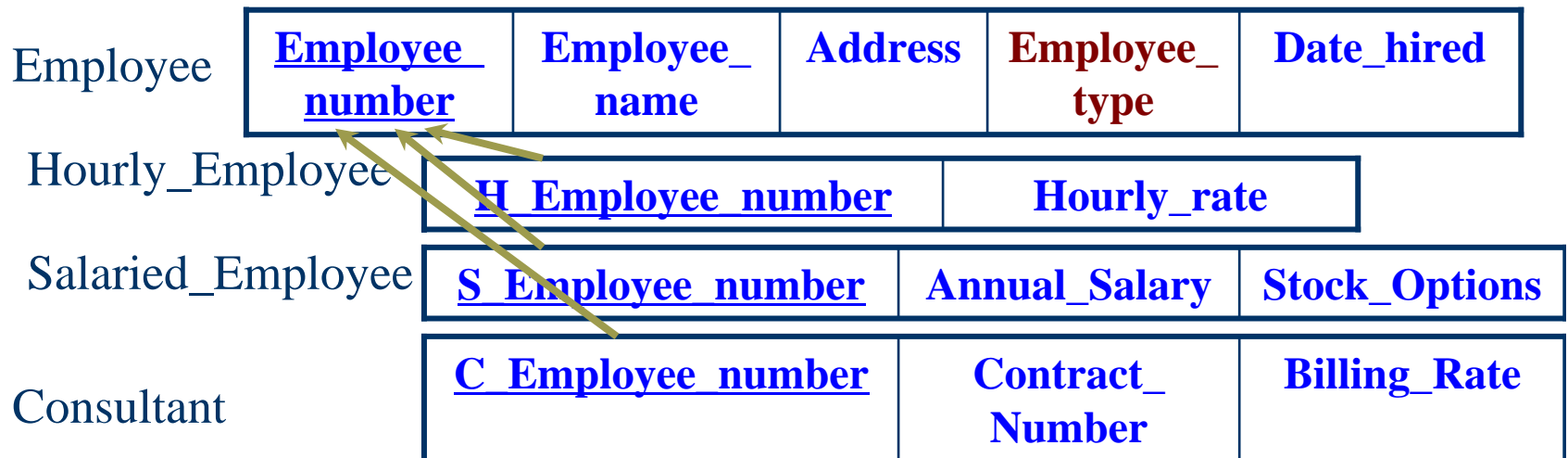
- ◆ Create one relation for each participating entity type
- ◆ Create a new associative relation:
  - The default primary key is the 3 primary keys of the 3 relations, also as the foreign keys respectively (maybe there are some additional ones)



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# Step7: Map supertype/subtype relationships (1)

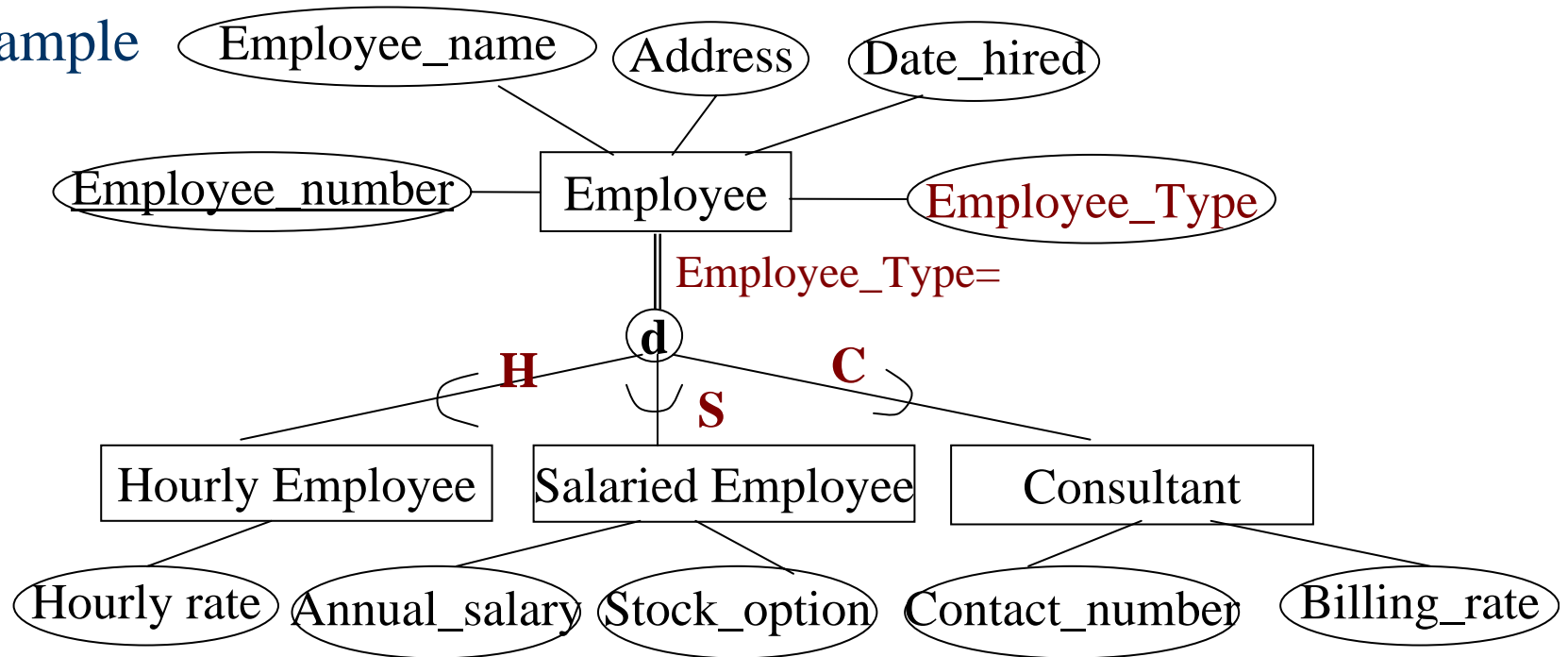
- ◆ Create a separate relation for the supertype and each of its subtypes, including the corresponding attributes
- ◆ Assign to the relation for each subtype the primary key of the supertype (also as the foreign keys respectively)
- ◆ Assign one or more attributes of the supertype as the subtype discriminator



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# Step 7: Map supertype/subtype relationships (2)

Example



How to get all the information for “Salaried Employee”?

# Introduction to normalization (The theory of database design)

- ◆ Motivation
- ◆ Preliminaries
- ◆ The basic normal forms

## **Motivation:**

- To get the well-structured relations
- To avoid unnecessary duplication and redundancy of data
- Decompose the relations and give the formal process
- ...

**Normalization:** the process of decomposing relations *with anomalies* to produce smaller, well-structured relations



# Preliminaries and basic concepts (1)

## ◆ Normal forms

First normal form, second normal form, third normal form, BC normal form, ...

## ◆ Functional dependencies(FDs)

- A constraint between two attributes or two sets of attributes

-  $A \rightarrow B$

① The value of A uniquely determines the value of B.

② Attribute B is functionally dependent on attribute A.

③ If the values on A are equal, then the values on B must be equal

E.g,  $Sno \rightarrow Sname$ ,  $Cno \rightarrow Cname$ ,  $Sno, Cno \rightarrow Grade$ ,

$ISBN \rightarrow Title, First\_Author\_Name$

③ The attribute may be functionally dependent on more than one attributes

④ Determinants: the attribute on the left-hand side of the arrow in a FD

# Preliminaries and basic concepts (2)

## ◆ Candidate keys and primary keys

- An attribute, or combination of attributes that **uniquely identifies** a row in a relation

① Unique identification

② Nonredundancy

E.g, Enroll(Sno, Cno, Grade)

◇ Sno → Grade? Cno → Grade? Sno, Cno → Grade?

◇ Either Sno or Cno can not respectively identify an instance in *Enroll* uniquely——The minimal attribute set as identifier

E.g, Shipment(Shipment\_no, Customer\_id, Vendor\_id, Date, Amount)

◇ Shipment\_no → Date, Amount Customer\_id, Vendor\_id → Date, Amount

◇ Both Shipment\_no and (Customer\_id, Vendor\_id) can be the unique identifier——Candidate keys☺!

# Preliminaries and basic concepts (3)

## ◆ **Primary key (PK)**

- The chosen candidate key (only one)
- Multiple candidate keys, unique primary key

E.g, Either **Shipment\_no** or (**Customer\_id, Vendor\_id**) be chosen.

## ◆ **Nonkey attributes**

Attributes except the key attributes

E.g, **Shipment(Shipment\_no, Customer\_id, Vendor\_id, Date, Amount)**

◇ nonkey attributes? **Date, Amount**

## ◆ **Candidate key revisited**

uniquely identifies the nonkey attributes (**all?**) in a relation

# Preliminaries and basic concepts (4)

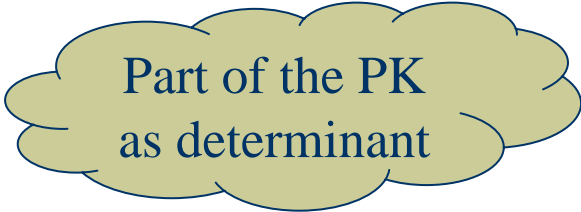
- ◆ **Some concepts of FDs**

E.g,  $\text{Employee2}(\underline{\text{Emp\_id}}, \text{Name}, \text{Dept\_name}, \text{Salary}, \underline{\text{Course\_title}}, \text{Date\_completed})$

- ◇ Primary key:  $\text{Emp\_id}, \text{Course\_title}$

- ◇  $\text{Emp\_id} \rightarrow \text{Name}, \text{Dept\_name}, \text{Salary}$

- ◇  $\text{Emp\_id}, \text{Course\_title} \rightarrow \text{Date\_completed}$



Part of the PK  
as determinant

- **Partial functional dependencies**

One or more nonkey attributes are functionally dependent on part of the PK

- **Full functional dependencies**

All nonkey attributes are functional dependent on the PK

E.g,  $\text{Shipment\_no} \rightarrow \text{Date}, \text{Amount}$

# Preliminaries and basic concepts (5)

## ◆ Some concepts of FDs (Cont.)

- **Trivial functional dependencies:**  $X \rightarrow Y$ , and  $Y \subseteq X$

E.g,  $\text{Shipment\_no} \rightarrow \text{Shipment\_no}$ ,

$\text{Customer\_id, Vendor\_id} \rightarrow \text{Customer\_id}$

- **Nontrivial functional dependencies** (Generally)

- **Transitive functional dependencies**

E.g,  $\text{Sales}(\underline{\text{Cust\_id}}, \text{Name}, \text{Salesperson}, \text{Region})$

◇  $\text{Cust\_id} \rightarrow \text{Name}, \text{Salesperson}, \text{Region}$

◇  $\text{Salesperson} \rightarrow \text{Region}$

◇  $\text{Cust\_id} \rightarrow \text{Salesperson}, \text{Salesperson} \rightarrow \text{Region}$

Region transitively dependent on Cust\_id

Between two or more nonkey attributes, such as  $A \rightarrow B, B \rightarrow C$

# Preliminaries and basic concepts (6)

- ◆ What results will be produced due to the partial and transitive FDs?  
Insertion, deletion, modification anomalies☹

E.g, The insertion and deletion anomaly with partial FD

<u>Emp_id</u>	Name	Dept_name	Salary	<u>Course title</u>	Date_completed
100	Simpson	Marketing	48,000	SPSS	6/19/200X
100	Simpson	Marketing	48,000	Surveys	10/7/200X
150	Martin	Marketing	42,000	Java	8/12/200X
...					

- How to insert a new course “C++”?
- What will be resulted when a course is deleted?

# The basic normal forms (1)

- ◆ **First normal form**

The relation contains no multivalued attributes

- ◆ **Second normal form**

- First normal form

- Every nonkey attribute is **fully functionally dependent** on the primary key

E.g, Employee2(Emp\_id, Name, Dept\_name, Salary, Course\_title,  
Date\_completed)

- ◇ Emp\_id → Name, Dept\_name, Salary

- Emp\_id, Course\_title → Date\_completed

- ◇ Decomposition:

- Employee1(Emp\_id, Name, Dept\_name, Salary )

- Emp\_Course(Emp\_id, Course\_title, Date\_completed)

# The basic normal forms (2)

## ◆ Third normal form

- Second normal form  $3NF \subseteq 2NF \subseteq 1NF$
- No transitive FDs present

E.g, Sales(Cust\_id, Name, Salesperson, Region)

- ◇  $Cust\_id \rightarrow Name, Salesperson, Region$
- ◇  $Salesperson \rightarrow Region$
- ◇ Decomposition:

Sales1(Cust\_id, Name, Salesperson )

Sperson(Salesperson, Region)

E.g, Shipment(Snum, Origin, Destination, Distance)

- ◇  $Snum \rightarrow Origin, Destination, Distance$  Shipto(Snum, Origin, Destination)
- ◇  $Origin, Destination \rightarrow Distance$  Distances(Origin, Destination, Distance)



# Discussion

- ◆ The seven steps of transforming EER diagram into relations are the experiential rules of logical database design
- ◆ The normalization and normal forms are the backbones in database design theory
- ◆ The seven steps of mapping are consistent with the normalization theory
- ◆ We can decompose a universal relational schema into well-structured ones based on the normal form decomposition algorithms
- ◆ How to perform the formal decomposition process?

***FDs are the most important thing in normalization!***

**How to get the enough FDs on the given relations?**

# Merging relations (1)

## ◆ Prerequisite

- Not the realm of database design theory
- From the viewpoint of database system development and view, information integration

## ◆ Motivation

- The work of several subteams comes together
- integrating existing databases, **merging tables**
- new data requirements arise

## ◆ Example

- Combine(merge) the relations that have the same PK  
Employee1(Employee\_id, Name, Address, Phone)  
Employee2(Employee\_id, Name, Address, Jobcode, No\_Years)
- Result of merging:  
Employee2(Employee\_id, Name, Address, Phone, Jobcode, No\_Years)

# Merging relations (2)

## ◆ View integration problems

### - Synonyms

Two or more attributes describing the same characteristics of an entity:

**Different names but the same meaning**

E.g, Student1(Student\_id, Name), Student2(Student\_no, Name, Address)

Merging result: Student(SSN, Name, Address)

### - Homonyms

An attribute that may have more than one meaning in various contexts

E.g, Student1(Student\_id, Name, Address),

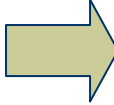
Student2(Student\_id, Name, Phone\_no, Address)

Merging result with specific prefixes:

Student(Student\_id, Name, Campus\_address, Permanent\_address)

# Merging relations (3)

- ◆ **Transitive FDs after the merging**

E.g,  $\text{Student1}(\underline{\text{Student\_id}}, \text{Major})$ ,  
 $\text{Student2}(\underline{\text{Student\_id}}, \text{Advisor})$    
 $\text{Student}(\underline{\text{Student\_id}}, \text{Major}, \text{Advisor})$

- ◇  $\text{Student\_id} \rightarrow \text{Major}, \text{Major} \rightarrow \text{Advisor}$
- ◇ Transitive FD is resulted
- ◇ Decomposition after the merging  
 $\text{Student}(\underline{\text{Student\_id}}, \text{Major})$   
 $\text{Major Advisor}(\underline{\text{Major}}, \text{Advisor})$

# Some other principles in the logical database design

## ◆ **Key definition**

- A primary key whose value is unique across all relations
- Similar to the object identifier
- PK maybe has no business meanings

## ◆ **Attribute definition**

- Some extra attributes are necessary in case of database evolution
- ...

## ◆ **The normalization is effectively used to check the transformation results from EER diagrams, and some complementary revision is doomed**

# Summary

- ◆ The relational data model
  - Definitions
  - Constraints: domain, entity, referential integrity (PK and FK)
  - Well-structured relation, anomalies
- ◆ Major steps in the logical database design based on transforming EER diagrams to normalized relations
  - Transform EER diagrams to relations (7 steps)
  - Normalize the relations (FDs, normal forms, schema decomposition)
  - Merge the relations (Synonyms, Homonyms)

*Spend more time on  
this chapter 😊!*

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# Assignments (1)

- ◆ **Define and contrast the following terms:**
  - (1) Candidate key; Primary key; Foreign key
  - (2) Entity integrity constraint; Referential integrity constraint
  - (3) Partial functional dependency; Transitive functional dependency
- ◆ **Describe the following concepts, as well as related practical examples (motivation, concept, and application):**
  - (1) Anomalies that arise in a table
  - (2) Foreign key, and referential integrity constraint
  - (3) Normalization and relational schema decomposition



# Assignments (2)

- ◆ Page 203, Review questions 17
- ◆ Page 203, Problems and exercises 3: (c), (d)
- ◆ Page 204, Problems and exercise 5: (a)
- ◆ Page 204, Problems and exercises 6: (a), (b), (c)
- ◆ Page 204, Problems and exercises 7: (a), (b), (c), (d)



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The end

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*Thanks!*



Logical database design and the  
relational model